

Input-Sensitive Profiling (or how to find the big-Oh of a program?)

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Conventional profilers collect cumulative data over a whole execution...



A possible approach is to extract and isolate the interesting code and perform multiple under a traditional profiler with different input but...



often hard to isolate portions of code and analyze them separately...

Input-Sensitive Profiling: aggregate routine times by input sizes

For routine f, collect a set of tuples, where each tuple contains:



No information about how performance of **single portions** of code **scales** as a function of the **input size**





typical usage scenarios...

RMS(f) = 2

RMS(g) = 3

Hard to collect real data about

Miss cache effects due to the interaction with the overall application...



- an estimate of an input size
- number of invocations on this input size
- max/min/avg execution cost

We need a metric for estimating the input size of a routine invocation...



How can measure the input size of a routine invocation **automatically**?



Read Memory Size: number of distinct memory cells first accessed by a routine, or by a descendent in the call tree, with a read operation

How can we compute **efficiently** the read memory size?

Two data structures:



- 1) a **shadow runtime stack**, where each entry contains:
- ID of pending routine
- routine entry timestamp
- total routine invocation cost
- partial read memory size



- more efficient/compact equal to the RMS upon invocation completion

2) a **shadow memory**:

t_{x}





procedure ca	ll(r):
top++	
S[top].rtn	←r

Case study: wf

We discuss **wf**, a simple word frequency counter included in the current development head of Fedora

Y Z

location w, each memory For timestamp ts[w] contains the time of **latest** access (read or write) to w

> aprof input-sensitive profiler based on



Comparable performance wrt other Valgrind tools. Experiments on CPU SPEC 2006 suite:



```
~30x
slowdown:
space overhead:
               ~2x
```

Profile data generated by aprof from a **single run** would require multiple runs of gprof

```
S[top].ts \leftarrow ++count
S[top].rms \leftarrow 0
S[top].cost \leftarrow get cost()
```

procedure return(): collect(S[top].rtn, S[top].rms, get cost() -S[top].cost) S[top - 1].rms += S[top].rms top----

```
procedure read(w):
  if ts[w ] < S[top].ts then</pre>
    S[top].rms++
    if ts[w] = 0 then
      let i be the max index in S
      such that S[i].ts \leq ts[w]
      S[i].rms—
    end if
  end if
  ts[w] \leftarrow count
```

16 256 >4000

procedure write(w): $ts[w] \leftarrow count$

Linux.

Our goal: study how the perfomance of individual routines scales as a function of the input size. To do so, for each routine of wf, we plot a chart with k points.

We analyze wf with:

	gprof	aprof
	For each point of a chart we need to perform a separate run of wf.	aprof can collects several points for a chart from the same execution of a pro- gram by aggregating rou- tine times by input sizes
	1 run = 1 point	1 run = N points
	Input of wf: texts of increasing size from classical literature	Input of wf: smallest text used with gprof
	Chart for str_tolower	Chart for str_tolower
]		

Profiles of CPU SPEC 2006 benchmarks: examples



Linear growth vs quadratic growth which one is correct? strlen() void str tolower(char* str) { int i; redundantly for (i = 0; i < strlen(str); i++)</pre> called at each str[i] = wf tolower(str[i]); iteration: $O(n^2)$ Fix the code by loop-invariant code motion: void str tolower(char* str) { Performance int i, len = strlen(str); improvement for (i = 0; i < len; i++)</pre> of wf up to str[i] = wf tolower(str[i]); 30% **Lesson**: input of str tolower are single words, not

the entire text. aprof automatically measures cost for each distinct word length.