

# Forward-Secure Sequential Aggregate Authentication



Di Ma and Gene Tsudik

Computer Science Department  
University of California, Irvine

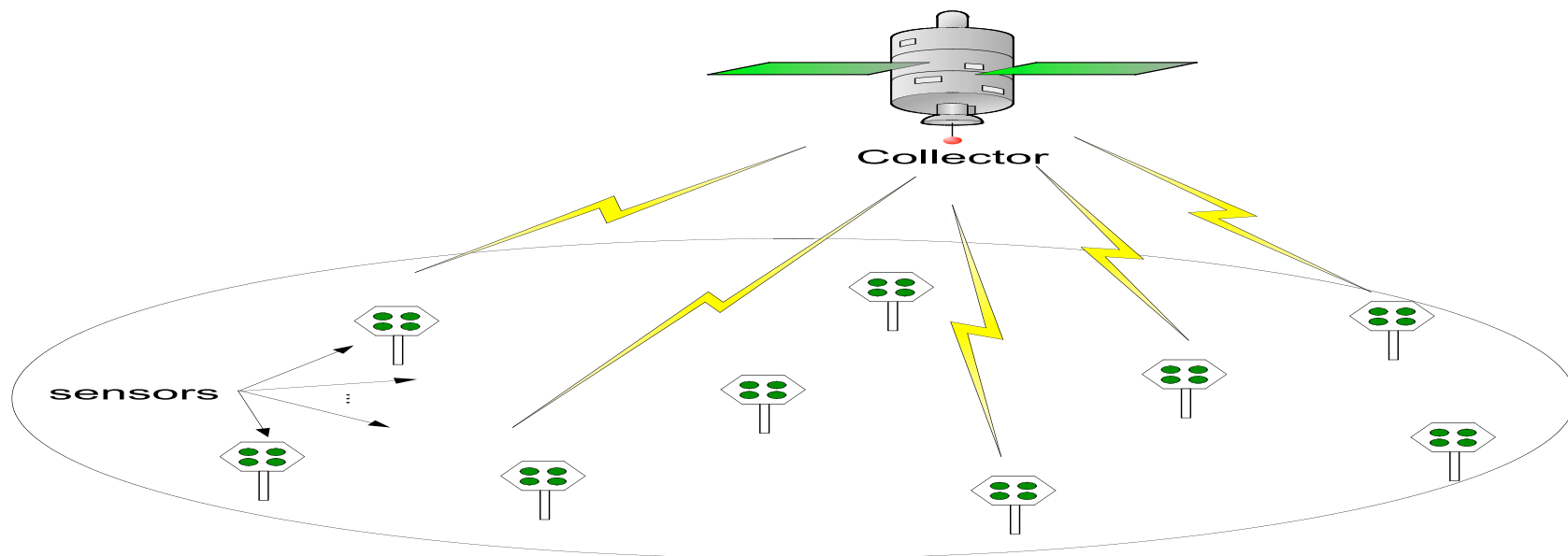
(Extended Abstract, presented at IEEE S&P 2007)

# Motivation

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## Unattended sensors:

- Sensors do not network
- Sensors unable to communicate to sink at will (in real time)
- Collect data and wait for collector
- Collector not fully trusted



## Goal:

authenticate data accumulated during multiple intervals

# Two Issues

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- Issue 1: Threat of Sensor Compromise
  - Sensor characteristics:
    - Low cost
    - No tamper-resistance
  - Adversarial and unattended environment
  - Compromise Possible --> need Forward Security:
    - Protect pre-compromise data
    - Periodic key evolution
- Issue 2: Storage and Communication Overhead
  - Sensors have limited on-board storage
  - Forward-security requires one authentication tag per message

# Forward-Secure Sequential Aggregate (*FssAgg*) Authentication

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- Reconcile minimal storage with mitigating potential sensor (and its signing key) compromise
- Allow **signer** to combine **multiple authentication tags** (generated within different time periods) into a **single constant-size tag**
  - Compromise of current key does not endanger authenticity of pre-compromise data
  - Verification of aggregate simultaneously verifies every component signature

# Definition

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## Component algorithms:

1. **FssAgg.Kg**: key generation, generate the initial signing key  $K_0$  and the verification key VK
2. **FssAgg.Asig**: aggregate sign, generate a FssAgg signature  $\sigma_{1,i}$  on message  $m_i$  and aggregate-so-far FssAgg signature  $\sigma_{1,i-1}$  with  $k_i$
3. **FssAgg.Upd**: key update, generate  $K_i$  from  $K_{i-1}$ , must be a one way function, and securely erase  $K_{i-1}$
4. **FssAgg.Aver**: aggregate verify, verify a FssAgg signature with the verification key VK, accept or reject

# Properties

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1. **Correctness:** any FssAgg signature produced by  $A_{sig}$  must be accepted by  $A_{ver}$ .
2. **Unforgeability:** without the knowledge of any signing keys, no adversary can compute an FssAgg signature
3. **Forward-security:** No adversary who breaks in  $i$ -th time period can generate a valid signature containing a signed message for any period  $j < i$

# MAC-based Scheme

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1. **FssAgg.Kg**: any symmetric key generation algorithm to generate  $k$ -bit secret  $s$  and set  $K_0 = VK = s$
2. **FssAgg.Asig**: for new message  $m_i$ :
  - a)  $\sigma_i = MAC(K_i, m_i)$
  - b)  $\sigma_{1,i} = H(\sigma_{1,i-1} \| \sigma_i)$
3. **FssAgg.Upd**:
  - a)  $K_{i+1} = H(K_i)$
  - b) Remove  $K_i$ , move to time  $i+1$
4. **FssAgg.Aver**: To verify  $\sigma_{1,i}$ :
  - a) Re-compute (from VK ):  $K_1 \dots K_i$ ,
  - b) Re-compute  $\sigma_{1,i}^c$
  - c) Compare  $\sigma_{1,i}^c \stackrel{?}{=} \sigma_{1,i}$

# MAC-based Scheme

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- Fast and space-efficient
- But:
  - Either collector cannot authenticate tags  
or
  - Collector can cheat
- Two MAC-based aggregates?
  - one for collector and one for sink
  - or signatures...



# Signature-based Scheme

- Boneh/Lynn/Shacham. Asiacrypt 2001
- Boneh/Gentry/Lynn/Shacham, Eurocrypt 2003

Based on BLS/BGLS signature scheme; works on groups with bilinear map  $e: G_1 \times G_2 \rightarrow G_T$  where :

- 1)  $G_1$  and  $G_2$  groups of order  $p$ ; 2)  $|G_1| = |G_2| = |G_T|$ ;
- 3)  $g_1, g_2$ : generator of  $G_1, G_2$

**1. FssAgg.Kg**: pick random  $x_0$  from  $Z_p$ , Compute pairs  $(x_i, v_i)$  s.t.  $x_i = H(x_{i-1})$ , and  $v_i = g_1^{x_i}$ , set  $K_0 = x_0$ ,  $VK = \{v_i\}$

**2. FssAgg.Asig**: given: new message  $m_i$ , current key  $K_i$ , and aggregate-so-far  $\sigma_{1,i-1}$

$$a) \sigma_i = BLS.sign(k_i, m_i)$$

$$b) \sigma_{1,i} = \sigma_{1,i-1} \bullet \sigma_i$$

**3. FssAgg.Update**:  $K_{i+1} = H(K_i)$ , remove  $K_i$

**4. FssAgg.Aver**: To verify  $\sigma_{1,i}$

$$e(\sigma_{1,i}, g_2) \stackrel{?}{=} \prod_{t=1}^i e(h_t, v_t)$$

# Performance Metrics

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## Signer efficiency

- Size of aggregate signature
  - Size of signing key
  - Complexity of key update
  - Complexity of aggregate signing
- Space
  - Time

## Verifier efficiency

- Size of verification key
- Complexity of aggregate verification

# Performance Evaluation

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- MAC scheme near-optimal
- Signature scheme is not
  - Signer-friendly
    - Constant private key and signature size
    - Efficient signing and key update
  - Not verifier-friendly
    - Public key size -  $O(T)$
    - Costly pairing operations in verification

# Summary and Future Work

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- Motivated Forward-Secure Aggregate authentication
- Two practical schemes:
  - MAC-based scheme near-optimal
  - Signature-based scheme not (yet) verifier-friendly
- **Future work:** more efficient schemes!

# Thanks!

